

# FLASH MEMORY

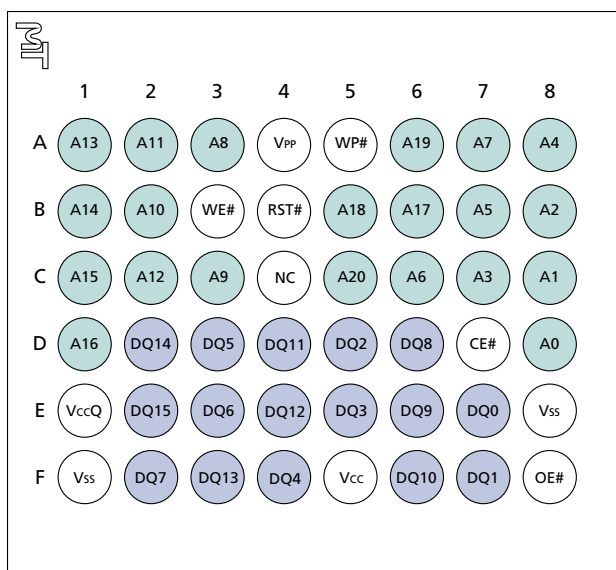
# MT28F322P3

**Low Voltage, Extended Temperature**

## FEATURES

- Flexible dual-bank architecture
  - Support for true concurrent operation with zero latency
  - Read bank *a* during program bank *b* and vice versa
  - Read bank *a* during erase bank *b* and vice versa
- Basic configuration:
  - Seventy-one erasable blocks
  - Bank *a* (8Mb for data storage)
  - Bank *b* (24Mb for program storage)
- V<sub>CC</sub>, V<sub>CCQ</sub>, V<sub>PP</sub> voltages
  - 2.7V (MIN), 3.3V (MAX) V<sub>CC</sub>
  - 2.2V (MIN), 3.3V (MAX) V<sub>CCQ</sub>
  - 3.0V (TYP) V<sub>PP</sub> (in-system PROGRAM/ERASE)
  - 12V ±5% (HV) V<sub>PP</sub> tolerant (factory programming compatibility)
- Random access time: 70ns @ 2.7V V<sub>CC</sub>
- Page Mode read access
  - Eight-word page
  - Interpage read access: 70ns @ 2.7V
  - Intrapage read access: 30ns @ 2.7V
- Low power consumption (V<sub>CC</sub> = 3.3V)
  - Asynchronous/interpage READ < 15mA
  - Intrapage READ < 7mA
  - WRITE < 20mA (MAX)
  - ERASE < 25mA (MAX)
  - Standby < 15µA (TYP), 50µA (MAX) @ 3.3V
  - Automatic power save (APS) feature
- Enhanced write and erase suspend options
  - ERASE-SUSPEND-to-READ within same bank
  - PROGRAM-SUSPEND-to-READ within same bank
  - ERASE-SUSPEND-to-PROGRAM within same bank
- Dual 64-bit chip protection registers for security purposes
- Cross-compatible command support
  - Extended command set
  - Common flash interface
- PROGRAM/ERASE cycle
  - 100,000 WRITE/ERASE cycles per block
- Fast programming algorithm
  - V<sub>PP</sub> = 12V ±5%

## BALL ASSIGNMENT 48-Ball FBGA



Top View  
(Ball Down)

**NOTE:** See page 7 for Ball Description Table.  
See page 35 for mechanical drawing.

## OPTIONS

- Timing
  - 70ns access
  - 80ns access
- Boot Block Configuration
  - Top
  - Bottom
- Package
  - 48-ball FBGA (6 x 8 ball grid)
- Operating Temperature Range
  - Commercial (0°C to +70°C)
  - Extended (-40°C to +85°C)

## MARKING

-70  
-80  
T  
B  
FJ  
None  
ET

Part Number Example:

**MT28F322P3FJ-70 BET**



## GENERAL DESCRIPTION

The MT28F322P3 is a high-performance, high-density, nonvolatile memory solution that can significantly improve system performance. This new architecture features a two-memory-bank configuration that supports background operation with no latency.

A high-performance bus interface allows a fast page mode data transfer; a conventional asynchronous bus interface is provided as well.

The MT28F322P3 allows soft protection for blocks, as read only, by configuring soft protection registers with dedicated command sequences. For security purposes, two 64-bit chip protection registers are provided.

The embedded WORD WRITE and BLOCK ERASE functions are fully automated by an on-chip write state machine (WSM). Two on-chip status registers, one for each of the two memory partitions, can be used to monitor the WSM status and to determine the progress of the program/erase task.

The erase/program suspend functionality allows compatibility with existing EEPROM emulation software packages.

The device is manufactured using 0.18µm process technology.

Please refer to Micron's Web site ([www.micron.com/flash](http://www.micron.com/flash)) for the latest data sheet.

## ARCHITECTURE AND MEMORY ORGANIZATION

The MT28F322P3 Flash device contains two separate banks of memory (bank *a* and bank *b*) for simultaneous READ and WRITE operations.

The MT28F322P3 Flash memory is available in the following bank segmentation configuration:

- Bank *a* comprises one-fourth of the memory and contains 8 x 4K-word parameter blocks and 15 x 32K-word blocks.
- Bank *b* represents three-fourths of the memory, is equally sectored, and contains 48 x 32K-word blocks.

Figures 2 and 3 show the bottom and top memory organizations.

## DEVICE MARKING

Due to the size of the package, Micron's standard part number is not printed on the top of each device. Instead, an abbreviated device mark comprised of a five-digit alphanumeric code is used. The abbreviated device marks are cross referenced to Micron part numbers in Table 1.

**Table 1**  
**Cross Reference for Abbreviated Device Marks**

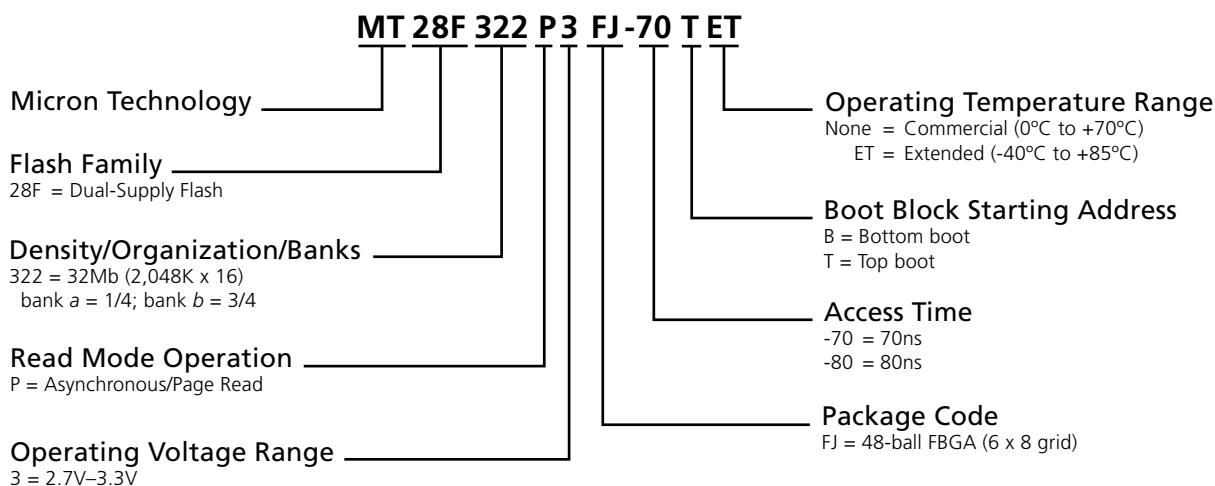
PART NUMBER	PRODUCT MARKING	SAMPLE MARKING	MECHANICAL SAMPLE MARKING
MT28F322P3FJ-70 BET	FW816	FX816	FY816
MT28F322P3FJ-70 TET	FW817	FX817	FY817
MT28F322P3FJ-80 BET	FW814	FX814	FY814
MT28F322P3FJ-80 TET	FW815	FX815	FY815

## PART NUMBERING INFORMATION

Micron's low-power devices are available with several different combinations of features (see Figure 1).

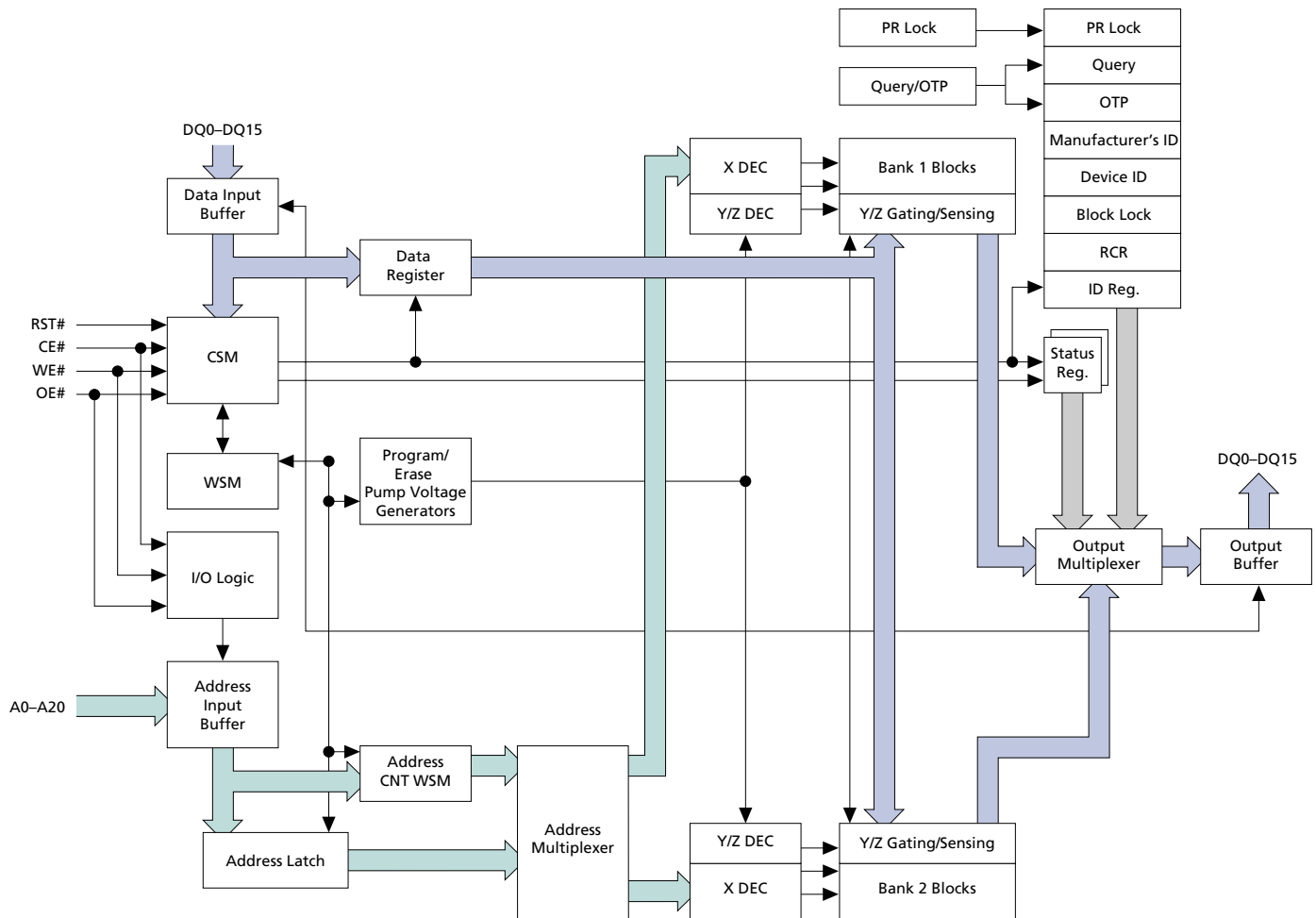
Valid combinations of features and their corresponding part numbers are listed in Table 2.

**Figure 1  
Part Number Chart**



**Table 2  
Valid Part Number Combinations**

PART NUMBER	ACCESS TIME (ns)	BOOT BLOCK STARTING ADDRESS	OPERATING TEMPERATURE RANGE
MT28F322P3FJ-70 BET	70	Bottom	-40°C to +85°C
MT28F322P3FJ-70 TET	70	Top	-40°C to +85°C
MT28F322P3FJ-80 BET	80	Bottom	-40°C to +85°C
MT28F322P3FJ-80 TET	80	Top	-40°C to +85°C

**FUNCTIONAL BLOCK DIAGRAM**




**Figure 2**  
**Bottom Boot Block Device**

Bank b = 24Mb		
Block	Block Size (K-bytes/K-words)	Address Range (x16)
70	64/32	1F8000h-1FFFFFh
69	64/32	1F0000h-1F7FFFh
68	64/32	1E8000h-1EFFFFh
67	64/32	1E0000h-1E7FFFh
66	64/32	1D8000h-1DFFFFh
65	64/32	1D0000h-1D7FFFh
64	64/32	1C8000h-1CFFFFh
63	64/32	1C0000h-1C7FFFh
62	64/32	1B8000h-1BFFFFh
61	64/32	1B0000h-1B7FFFh
60	64/32	1A8000h-1AFFFFh
59	64/32	1A0000h-1A7FFFh
58	64/32	198000h-19FFFFh
57	64/32	190000h-197FFFh
56	64/32	188000h-18FFFFh
55	64/32	180000h-187FFFh
54	64/32	178000h-17FFFFh
53	64/32	170000h-177FFFh
52	64/32	168000h-16FFFFh
51	64/32	160000h-167FFFh
50	64/32	158000h-15FFFFh
49	64/32	150000h-157FFFh
48	64/32	148000h-14FFFFh
47	64/32	140000h-147FFFh
46	64/32	138000h-13FFFFh
45	64/32	130000h-137FFFh
44	64/32	128000h-12FFFFh
43	64/32	120000h-127FFFh
42	64/32	118000h-11FFFFh
41	64/32	110000h-117FFFh
40	64/32	108000h-10FFFFh
39	64/32	100000h-107FFFh
38	64/32	0F8000h-0FFFFFh
37	64/32	0F0000h-0F7FFFh
36	64/32	0E8000h-0EFFFFh
35	64/32	0E0000h-0E7FFFh
34	64/32	0D8000h-0DFFFFh
33	64/32	0D0000h-0D7FFFh
32	64/32	0C8000h-0CFFFFh
31	64/32	0C0000h-0C7FFFh
30	64/32	0B8000h-0BFFFFh
29	64/32	0B0000h-0B7FFFh
28	64/32	0A8000h-0AFFFFh
27	64/32	0A0000h-0A7FFFh
26	64/32	098000h-097FFFh
25	64/32	090000h-097FFFh
24	64/32	088000h-087FFFh
23	64/32	080000h-087FFFh

Bank a = 8Mb		
Block	Block Size (K-bytes/K-words)	Address Range (x16)
22	64/32	078000h-07FFFFh
21	64/32	070000h-077FFFh
20	64/32	068000h-067FFFh
19	64/32	060000h-067FFFh
18	64/32	058000h-05FFFFh
17	64/32	050000h-057FFFh
16	64/32	048000h-04FFFFh
15	64/32	040000h-047FFFh
14	64/32	038000h-03FFFFh
13	64/32	030000h-037FFFh
12	64/32	028000h-02FFFFh
11	64/32	020000h-027FFFh
10	64/32	018000h-01FFFFh
9	64/32	010000h-017FFFh
8	64/32	008000h-00FFFFh
7	8/4	007000h-007FFFh
6	8/4	006000h-006FFFh
5	8/4	005000h-005FFFh
4	8/4	004000h-004FFFh
3	8/4	003000h-003FFFh
2	8/4	002000h-002FFFh
1	8/4	001000h-001FFFh
0	8/4	000000h-000FFFh

**Figure 3  
Top Boot Block Device**

Bank a = 8Mb		
Block	Block Size (K-bytes/K-words)	Address Range (x16)
70	8/4	1FF000h-1FFFFFh
69	8/4	1FE000h-1FEFFFh
68	8/4	1FD000h-1FDFFFh
67	8/4	1FC000h-1FCFFFh
66	8/4	1FB000h-1FBFFFh
65	8/4	1FA000h-1FAFFFh
64	8/4	1F9000h-1F9FFFh
63	8/4	1F8000h-1F8FFFh
62	64/32	1F0000h-1F7FFFh
61	64/32	1E8000h-1EFFFFh
60	64/32	1E0000h-1E7FFFh
59	64/32	1D8000h-1DFFFFh
58	64/32	1D0000h-1D7FFFh
57	64/32	1C8000h-1C7FFFh
56	64/32	1C0000h-1C7FFFh
55	64/32	1B8000h-1BFFFFh
54	64/32	1B0000h-1B7FFFh
53	64/32	1A8000h-1AFFFFh
52	64/32	1A0000h-1A7FFFh
51	64/32	198000h-19FFFFh
50	64/32	190000h-197FFFh
49	64/32	188000h-18FFFFh
48	64/32	180000h-187FFFh

Bank b = 24Mb		
Block	Block Size (K-bytes/K-words)	Address Range (x16)
47	64/32	178000h-17FFFFh
46	64/32	170000h-177FFFh
45	64/32	168000h-16FFFFh
44	64/32	160000h-167FFFh
43	64/32	158000h-15FFFFh
42	64/32	150000h-157FFFh
41	64/32	148000h-14FFFFh
40	64/32	140000h-147FFFh
39	64/32	138000h-13FFFFh
38	64/32	130000h-137FFFh
37	64/32	128000h-12FFFFh
36	64/32	120000h-127FFFh
35	64/32	118000h-11FFFFh
34	64/32	110000h-117FFFh
33	64/32	108000h-10FFFFh
32	64/32	100000h-107FFFh
31	64/32	0F8000h-0FFFFFh
30	64/32	0F0000h-0F7FFFh
29	64/32	0E8000h-0EFFFFh
28	64/32	0E0000h-0E7FFFh
27	64/32	0D8000h-0DFFFFh
26	64/32	0D0000h-0D7FFFh
25	64/32	0C8000h-0CFFFFh
24	64/32	0C0000h-0C7FFFh
23	64/32	0B8000h-0BFFFFh
22	64/32	0B0000h-0B7FFFh
21	64/32	0A8000h-0AFFFFh
20	64/32	0A0000h-0A7FFFh
19	64/32	098000h-09FFFFh
18	64/32	090000h-097FFFh
17	64/32	088000h-08FFFFh
16	64/32	080000h-087FFFh
15	64/32	078000h-07FFFFh
14	64/32	070000h-077FFFh
13	64/32	068000h-06FFFFh
12	64/32	060000h-067FFFh
11	64/32	058000h-05FFFFh
10	64/32	050000h-057FFFh
9	64/32	048000h-04FFFFh
8	64/32	040000h-047FFFh
7	64/32	038000h-03FFFFh
6	64/32	030000h-037FFFh
5	64/32	028000h-02FFFFh
4	64/32	020000h-027FFFh
3	64/32	018000h-01FFFFh
2	64/32	010000h-017FFFh
1	64/32	008000h-00FFFFh
0	64/32	000000h-007FFFh

**BALL DESCRIPTIONS**

48-BALL FBGA NUMBERS	SYMBOL	TYPE	DESCRIPTION
D8, C8, B8, C7, A8, B7, C6, A7, A3, C3, B2, A2, C2, A1, B1, C1, D1, B6, B5, A6, C5	A0–A20	Input	Address Inputs: Inputs for the address during READ and WRITE operations. Addresses are internally latched during READ and WRITE cycles.
D7	CE#	Input	Chip Enable: Activates the device when LOW. When CE# is HIGH, the device is disabled and goes into standby power mode.
F8	OE#	Input	Output Enable: Enables the output buffer when LOW. When OE# is HIGH, the output buffers are disabled.
B3	WE#	Input	Write Enable: Determines if a given cycle is a WRITE cycle. If WE# is LOW, the cycle is either a WRITE to the command state machine (CSM) or to the memory array.
B4	RST#	Input	Reset: When RST# is a logic LOW, the device is in reset mode, which drives the outputs to High-Z and resets the write state machine (WSM). When RST# is at logic HIGH, the device is in standard operation. When RST# transitions from logic LOW to logic HIGH, the device resets all blocks to locked and defaults to the read array mode.
A5	WP#	Input	Write Protect: Controls the lock down function of the flexible locking feature.
A4	V <sub>PP</sub>	Input	Program/Erase Enable: [1.8V–3.3V] Operates as input at logic levels to control complete device protection. Provides factory programming compatibility when driven to 11.4V–12.6V.
E7, F7, D5, E5, F4, D3, E3, F2, D6, E6, F6, D4, E4, F3, D2, E2	DQ0–DQ15	Input/ Output	Data Inputs/Outputs: Input array data on the second CE# and WE# cycle during PROGRAM command. Input commands to the command user interface when CE# and WE# are active. DQ0–DQ15 output data when CE# and OE# are active.
E8, F1	V <sub>SS</sub>	Supply	Do not float any ground ball.
F5	V <sub>CC</sub>	Supply	Device Power Supply: [2.7V–3.3V] Supplies power for device operation.
E1	V <sub>CCQ</sub>	Supply	I/O Power Supply: [2.2V–3.3V] Supplies power for input/output buffers.
C4	NC	–	Internally not connected.

## COMMAND STATE MACHINE (CSM)

Commands are issued to the command state machine (CSM) using standard microprocessor write timings. The CSM acts as an interface between external microprocessors and the internal WSM. The available commands are listed in Table 3, their definitions are given in Table 4, and their descriptions in Table 5. Program and erase algorithms are automated by an on-chip WSM. For more specific information about the CSM transition states, see Micron technical note TN-28-33, "Command State Machine Description and Command Definition."

Once a valid PROGRAM/ERASE command is entered, the WSM executes the appropriate algorithm, which generates the necessary timing signals to control the device internally and accomplish the requested operation. A command is valid only if the exact sequence of WRITES is completed. After the WSM completes its task, the WSM status bit (SR7) (see Table 7) is set to a logic HIGH level (1), allowing the CSM to respond to the full command set again.

## OPERATIONS

Device operations are selected by entering a standard JEDEC 8-bit command code with conventional microprocessor timings into an on-chip CSM through I/Os DQ0–DQ7. The number of bus cycles required to activate a command is typically one or two. The first operation is always a WRITE. Control signals CE# and WE# must be at a logic LOW level ( $V_{IL}$ ), and OE# and RST# must be at logic HIGH ( $V_{IH}$ ). The second operation, when needed, can be a WRITE or a READ depending upon the command. During a READ operation, con-

trol signals CE# and OE# must be at a logic LOW level ( $V_{IL}$ ), and WE# and RST# must be at logic HIGH ( $V_{IH}$ ).

Table 6 shows the bus operations for all the modes: write, read, reset, standby, and output disable.

When the device is powered up, internal reset circuitry initializes the chip to a read array mode of operation. Changing the mode of operation requires that a command code be entered into the CSM. For each one of the two memory partitions, an on-chip status register is available. These two registers allow the progress of the various operations that can take place on a memory bank to be monitored. One of the two status registers is interrogated by entering a READ STATUS REGISTER command onto the CSM (cycle 1), specifying an address within the memory partition boundary, and reading the register data on I/Os DQ0–DQ7 (cycle 2). Status register bits SR0–SR7 correspond to DQ0–DQ7 (see Table 7).

## COMMAND DEFINITION

Once a specific command code has been entered, the WSM executes an internal algorithm, generating the necessary timing signals to program, erase, and verify data. See Table 4 for the CSM command definitions and data for each of the bus cycles.

## STATUS REGISTER

The status register allows the user to determine whether the state of a PROGRAM/ERASE operation is pending or complete. The status register is monitored by toggling OE# and CE# and reading the resulting status code on I/Os DQ0–DQ7. The high-order I/Os

**Table 3**  
**Command State Machine Codes For Device Mode Selection**

COMMAND DQ0–DQ7	CODE ON DEVICE MODE
40h/10h	Program setup/alternate program setup
20h	Block erase setup
30h	Fast programming algorithm setup
50h	Clear status register
60h	Protection configuration setup
70h	Read status register
90h	Read protection configuration register
98h	Read query
B0h	Program/erase suspend
C0h	Protection register program/lock
D0h	Program/erase resume – erase confirm
FFh	Read array



(DQ8–DQ15) are set to 00h internally, so only the low-order I/Os (DQ0–DQ7) need to be interpreted. Address lines select the status register pertinent to the selected memory partition.

Register data is updated and latched on the falling edge of OE# or CE#, whichever occurs last. Latching the data prevents errors from occurring if the register input changes during a status register read.

The status register provides the internal state of the WSM to the external microprocessor. During periods when the WSM is active, the status register can be polled to determine the WSM status. Table 7 defines the status register bits.

After monitoring the status register during a PROGRAM/ERASE operation, the data appearing on DQ0–DQ7 remains as status register data until a new command is issued to the CSM. To return the device to other modes of operation, a new command must be issued to the CSM.

## CSM OPERATIONS

The CSM decodes instructions for read array, read protection configuration register, read query, read status register, clear status register, program, erase, erase suspend, erase resume, program suspend, program resume, lock block, unlock block, and lock down block, chip protection program, and set read configuration register. The 8-bit command code is input to the device on DQ0–DQ7 (see Table 3 for CSM codes and Table 4 for command definitions). During a PROGRAM or ERASE cycle, the CSM informs the WSM that a PROGRAM or ERASE cycle has been requested.

During a PROGRAM cycle, the WSM controls the program sequences and the CSM responds to a PROGRAM SUSPEND command only.

During an ERASE cycle, the CSM responds to an ERASE SUSPEND command only. When the WSM has completed its task, the WSM status bit (SR7) is set to a

**Table 4**  
**Command Definitions**

COMMAND	FIRST BUS CYCLE			SECOND BUS CYCLE		
	OPERATION	ADDRESS <sup>1</sup>	DATA	OPERATION	ADDRESS <sup>1</sup>	DATA <sup>1</sup>
READ ARRAY	WRITE	WA	FFh			
READ PROTECTION CONFIGURATION REGISTER	WRITE	IA	90h	READ	IA	ID
READ STATUS REGISTER	WRITE	BA	70h	READ	X	SRD
CLEAR STATUS REGISTER	WRITE	BA	50h			
READ QUERY	WRITE	QA	98h	READ	QA	QD
BLOCK ERASE SETUP	WRITE	BA	20h	WRITE	BA	D0h
PROGRAM SETUP/ALTERNATE PROGRAM SETUP	WRITE	WA	40h/10h	WRITE	WA	WD
FAST PROGRAMMING ALGORITHM SETUP	WRITE	WA	30h	WRITE	WA	D0h
PROGRAM/ERASE SUSPEND	WRITE	BA	B0h			
PROGRAM/ERASE RESUME - ERASE CONFIRM	WRITE	BA	D0h			
LOCK BLOCK	WRITE	BA	60h	WRITE	BA	01h
UNLOCK BLOCK	WRITE	BA	60h	WRITE	BA	D0h
LOCK DOWN BLOCK	WRITE	BA	60h	WRITE	BA	2Fh
PROTECTION REGISTER PROGRAM	WRITE	PA	C0h	WRITE	PA	PD
PROTECTION REGISTER LOCK	WRITE	LPA	C0h	WRITE	LPA	FFDh

**NOTE:** 1. BA: Address within the block  
IA: Identification code address  
ID: Identification code data  
LPA: Lock protection register address  
PA: Protection register address  
PD: Data to be written at location PA  
QA: Query code address

QD: Query code data  
SRD: Data read from the status register  
WA: Word address of memory location to be written, or read  
WD: Data to be written at the location WA  
X: "Don't Care"

**Table 5  
Command Descriptions**

CODE	DEVICE MODE	BUS CYCLE	DESCRIPTION
10h	Alt. Program Setup	First	Operates the same as a PROGRAM SETUP command.
20h	Erase Setup	First	Prepares the CSM for an ERASE CONFIRM command. If the next command is not an ERASE CONFIRM command, the command will be ignored, and the bank will go to the read array mode and wait for another command.
30h	FPA Setup	First	Prepares the CSM for an FPA CONFIRM command.
40h	Program Setup	First	A two-cycle command: The first cycle prepares for a PROGRAM operation, and the second cycle latches addresses and data and initiates the WSM to execute the program algorithm. The Flash device outputs status register data on the falling edge of OE# or CE#, whichever occurs first.
50h	Clear Status Register	First	The WSM can set the block lock status (SR3), program status (SR4), and erase status (SR5) bits in the status register to "1," but it cannot clear them to "0." Issuing this command clears those bits to "0."
60h	Protection Configuration Setup	First	Prepares the CSM for changes to the block locking status. If the next command is not BLOCK UNLOCK, BLOCK LOCK or BLOCK LOCK DOWN, the command will be ignored, and the device will go to read status mode.
70h	Read Status Register	First	Places the device into read status register mode. Reading the device will output the contents of the status register for the addressed bank. The device will automatically enter this mode for the addressed bank after a PROGRAM or ERASE operation has been initiated.
90h	Read Protection Configuration	First	Puts the device into the read protection configuration mode so that reading the device will output the manufacturer/device codes or block lock status.
98h	Read Query	First	Puts the device into the read query mode so that reading the device will output common Flash interface information.
B0h	Program/Erase Suspend	First	Suspends the currently executing PROGRAM/ERASE operation. The status register will indicate when the operation has been successfully suspended by setting either the program suspend (SR2) or erase suspend (SR6), and the WSM status bit (SR7) to a "1" (ready). The WSM will continue to idle in the suspend state, regardless of the state of all input control signals except RST#, which will immediately shut down the WSM and the remainder of the chip if RST# is driven to V <sub>IL</sub> .
C0h	Program Device Protection Register	First	Writes a specific code into the device protection register.
	Lock Device Protection Register	First	Locks the device protection register; data can no longer be changed.

(continued on the next page)

**Table 5**  
**Command Descriptions (continued)**

CODE	DEVICE MODE	BUS CYCLE	DESCRIPTION
D0h	Erase Confirm	Second	If the previous command was an ERASE SETUP command, then the CSM will close the address and data latches, and it will begin erasing the block indicated on the address bus. During programming/erase, the device will respond only to the READ STATUS REGISTER, PROGRAM SUSPEND, or ERASE SUSPEND commands and will output status register data on the falling edge of OE# or CE#, whichever occurs last.
	Program/Erase Resume	First	If a PROGRAM or ERASE operation was previously suspended, this command will resume the operation.
D0h	FPA Confirm	Second	If the previous command was FPA SETUP, the CSM will latch the address indicated on the address bus and enter the FPA mode.
FFh	Read Array	First	During the array mode, array data will be output on the data bus.
01h	Lock Block	Second	If the previous command was PROTECTION CONFIGURATION SETUP, the CSM will latch the address and lock the block indicated on the address bus.
2Fh	Lock Down	Second	If the previous command was PROTECTION CONFIGURATION SETUP, the CSM will latch the address and lock down the block indicated on the address bus.
D0h	Unlock Block	Second	If the previous command was PROTECTION CONFIGURATION SETUP, the CSM will latch the address and unlock the block indicated on the address bus. If the block had been previously set to lock down, this operation will have no effect.
00h	Invalid /Reserved		Unassigned command that should not be used.

logic HIGH level and the CSM responds to the full command set. The CSM stays in the current command state until the microprocessor issues another command.

The WSM successfully initiates an ERASE or PROGRAM operation only when  $V_{PP}$  is within its correct voltage range.

### **CLEAR STATUS REGISTER**

The internal circuitry can set, but not clear, the block lock status bit (SR1), the  $V_{PP}$  status bit (SR3), the program status bit (SR4), and the erase status bit (SR5) of the status register. The CLEAR STATUS REGISTER command (50h) allows the external microprocessor to clear these status bits and synchronize to the internal operations. When the status bits are cleared, the device returns to the read array mode.

### **READ OPERATIONS**

The following READ operations are available: READ ARRAY, READ PROTECTION CONFIGURATION REGISTER, READ QUERY and READ STATUS REGISTER.

### **READ ARRAY**

The array is read by entering the command code FFh on DQ0–DQ7. Control signals CE# and OE# must be at a logic LOW level ( $V_{IL}$ ), and WE# and RST# must be at logic HIGH level ( $V_{IH}$ ) to read data from the array. Data is available on DQ0–DQ15. Any valid address within any of the blocks selects that address and allows data to be read from that address. Upon initial power-up or device reset, the device defaults to the read array mode.

### **READ PROTECTION CONFIGURATION DATA**

The chip identification mode outputs three types of information: the manufacturer/device identifier, the block locking status, and the protection register. Two bus cycles are required for this operation: the chip identification data is read by entering the command code 90h on DQ0–DQ7 to the bank containing address 0h and the identification code address on the address lines. Control signals CE# and OE# must be at a logic LOW level ( $V_{IL}$ ), and WE# and RST# must be at a logic HIGH level ( $V_{IH}$ ) to read data from the protection con-

figuration register. Data is available on DQ0–DQ15. After data is read from the protection configuration register, the read array command, FFh, must be issued to the bank containing address 00h prior to issuing other commands. See Table 9 for further details.

### READ QUERY

The read query mode outputs common flash interface (CFI) data when the device is read (see Table 11). Two bus cycles are required for this operation. It is possible to access the query by writing the read query command code 98h on DQ0–DQ7 to the bank containing address 0h. Control signals CE# and OE# must be at a logic LOW level ( $V_{IL}$ ), and WE# and RST# must be at a logic HIGH level ( $V_{IH}$ ) to read data from the query. The CFI data structure contains information such as block size, density, command set, and electrical specifications. To return to read array mode, write the read array command code FFh on DQ0–DQ7.

### READ STATUS REGISTER

The status register is read by entering the command code 70h on DQ0–DQ7. Two bus cycles are required for this operation: one to enter the command code and the block address and a second to read the status register. In a READ cycle, the address is latched and register data is updated on the falling edge of OE# or CE#, whichever occurs last.

## PROGRAMMING OPERATIONS

There are two CSM commands for programming: PROGRAM SETUP and ALTERNATE PROGRAM SETUP (see Table 3).

After the desired command code is entered (10h or 40h command code on DQ0–DQ7), the WSM takes over and correctly sequences the device to complete the PROGRAM operation. The WRITE operation may be monitored through the status register (see the Status Register section). During this time, the CSM will only respond to a PROGRAM SUSPEND command until the PROGRAM operation has been completed, after which time all commands to the CSM become valid again. The PROGRAM operation can be suspended by issuing a PROGRAM SUSPEND command (B0h). Once the WSM reaches the suspend state, it allows the CSM to respond only to READ ARRAY, READ STATUS REGISTER, READ PROTECTION CONFIGURATION, READ QUERY, PROGRAM SETUP, or PROGRAM RESUME. During the PROGRAM SUSPEND operation, array data should be read from an address other than the one being programmed. To resume the PROGRAM operation, a PROGRAM RESUME command (D0h) must be issued to cause the CSM to clear the suspend state previously

set (see Figure 4 for programming operation and Figure 5 for program suspend and program resume).

Taking RST# to  $V_{IL}$  during programming aborts the PROGRAM operation. During programming,  $V_{PP}$  must remain in the appropriate  $V_{PP}$  voltage range as shown in the recommended operating conditions table.

### FAST PROGRAMMING ALGORITHM (FPA) MODE

The fast programming algorithm (FPA) is intended for in-factory use. It enables fast data stream programming.

For in-factory programming, the FPA, along with an optimized set of programming parameters, minimizes chip programming time when  $11.4V \leq V_{PP} \leq 12.6V$ .

Executing the FPA command (30h), followed by FPA CONFIRM (D0h), enables an entire block to be programmed. This eliminates the need to continuously update the address to be programmed.

An initial delay is required after issuing the FPA command. (See the Erase and Program Cycle Timing Requirements Table.) The delay enables the device to detect 12V on  $V_{PP}$ . If  $V_{PP} < 11.4V$ , or if the block is locked, the status register returns an error. When the FPA command is executed successfully, a data stream can be programmed beginning at the first address. The address can be held constant, or it can be incremented within the address range. The program ends when the programmer enters an address outside the address range of the current block.

When the FPA is activated, the data must be provided in sequential order to the WSM. Immediately after programming, verification is executed. The address sequence is again provided to the WSM, which automatically performs a data consistency check between the data stored in the memory array, and the programmed data. The result is stored in the status register. Issuing an address outside the memory block boundaries exits the verification cycle. Figure 8 shows the FPA flowchart.

### ERASE OPERATIONS

An ERASE operation must be used to initialize all bits in an array block to “1s.” After BLOCK ERASE confirm is issued, the CSM responds only to an ERASE SUSPEND command until the WSM completes its task.

Block erasure inside the memory array sets all bits within the address block to logic 1s. Erase is accomplished only by blocks; data at single address locations within the array cannot be erased individually. The block to be erased is selected by using any valid address within that block. Block erasure is initiated by a command sequence to the CSM: BLOCK ERASE setup

(20h) followed by BLOCK ERASE CONFIRM (D0h) (see Figure 6). A two-command erase sequence protects against accidental erasure of memory contents.

When the BLOCK ERASE CONFIRM command is complete, the WSM automatically executes a sequence of events to complete the block erasure. During this sequence, the block is programmed with logic 0s, data is verified, all bits in the block are erased, and finally verification is performed to ensure that all bits are correctly erased. The ERASE operation may be monitored through the status register (see the Status Register section).

During the execution of an ERASE operation the ERASE SUSPEND command (B0h) can be entered to direct the WSM to suspend the ERASE operation. Once

the WSM has reached the suspend state, it allows the CSM to respond only to the READ ARRAY, READ STATUS REGISTER, READ QUERY, READ CHIP PROTECTION CONFIGURATION, PROGRAM SETUP, PROGRAM RESUME, ERASE RESUME and LOCK SETUP (see the Block Locking section). During the ERASE SUSPEND operation, array data must be read from a block other than the one being erased. To resume the ERASE operation, an ERASE RESUME command (D0h) must be issued to cause the CSM to clear the suspend state previously set (see Figure 7). It is also possible to suspend an ERASE in any bank and initiate a WRITE to another block in the same bank. After the completion of a WRITE, an ERASE can be resumed by writing an ERASE RESUME command.

**Table 6**  
**Bus Operations**

MODE	RST#	CE#	OE#	WE#	ADDRESS	DQ0-DQ15
Read (array, status registers, device identification register, or query)	V <sub>IH</sub>	V <sub>IL</sub>	V <sub>IL</sub>	V <sub>IH</sub>	X	DOUT
Standby	V <sub>IH</sub>	V <sub>IH</sub>	X	X	X	High-Z
Output Disable	V <sub>IH</sub>	V <sub>IH</sub>	X	X	X	High-Z
Reset	V <sub>IL</sub>	X	X	X	X	High-Z
Write	V <sub>IH</sub>	V <sub>IL</sub>	V <sub>IH</sub>	V <sub>IL</sub>	X	DIN

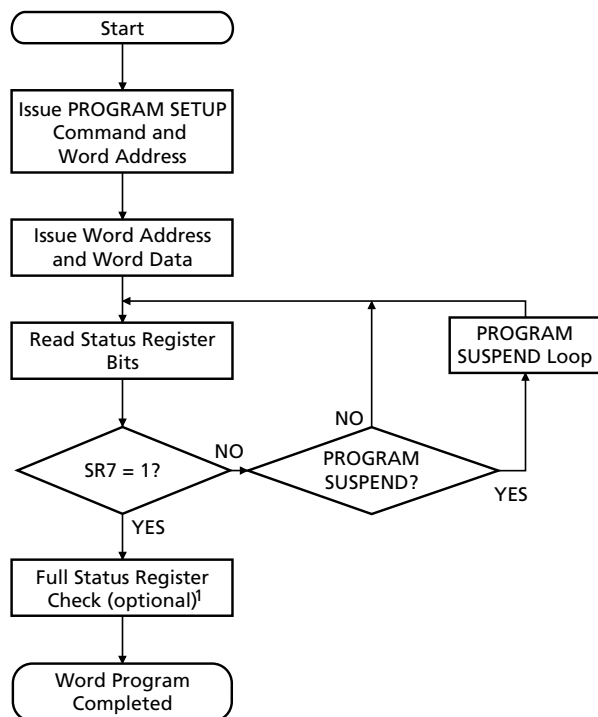
**Table 7**  
**Status Register Bit Definition**

<b>WSMS</b>	<b>ESS</b>	<b>ES</b>	<b>PS</b>	<b>V<sub>PP</sub>S</b>	<b>PSS</b>	<b>BLS</b>	<b>FPAS</b>
7	6	5	4	3	2	1	0

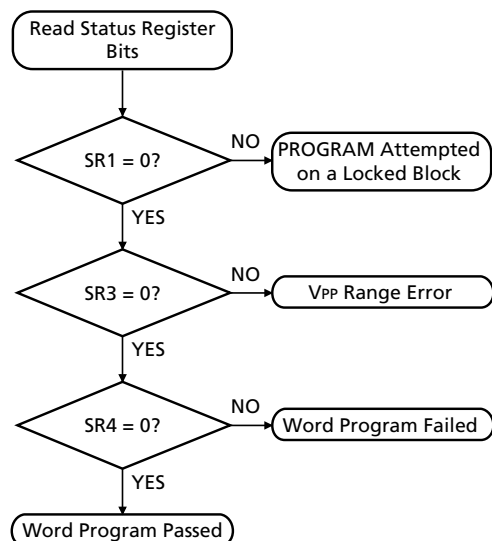
<b>STATUS BIT #</b>	<b>STATUS REGISTER BIT</b>	<b>DESCRIPTION</b>
SR7	WRITE STATE MACHINE STATUS (WSMS) 1 = Ready 0 = Busy	Check write state machine bit first to determine word program or block erase completion, before checking program or erase status bits.
SR6	ERASE SUSPEND STATUS (ESS) 1 = BLOCK ERASE Suspended 0 = BLOCK ERASE in Progress/Completed	When ERASE SUSPEND is issued, WSM halts execution and sets both WSMS and ESS bits to "1." ESS bit remains set to "1" until an ERASE RESUME command is issued.
SR5	ERASE STATUS (ES) 1 = Error in Block Erasure 0 = Successful BLOCK ERASE	When this bit is set to "1," WSM has applied the maximum number of erase pulses to the block and is still unable to verify successful block erasure.
SR4	PROGRAM STATUS (PS) 1 = Error in PROGRAM 0 = Successful PROGRAM	When this bit is set to "1," WSM has attempted but failed to program a word.
SR3	V <sub>PP</sub> STATUS (V <sub>PP</sub> S) 1 = V <sub>PP</sub> Low Detect, Operation Abort 0 = V <sub>PP</sub> = OK	The V <sub>PP</sub> status bit does not provide continuous indication of the V <sub>PP</sub> level. The WSM interrogates the V <sub>PP</sub> level only after the program or erase command sequences have been entered and informs the system if V <sub>PP</sub> < 1.8V. The V <sub>PP</sub> level is also checked before the PROGRAM/ERASE operation is verified by the WSM.
SR2	PROGRAM SUSPEND STATUS (PSS) 1 = PROGRAM Suspended 0 = PROGRAM in Progress/Completed	When PROGRAM SUSPEND is issued, WSM halts execution and sets both WSMS and PSS bits to "1." PSS bit remains set to "1" until a PROGRAM RESUME command is issued.
SR1	BLOCK LOCK STATUS (BLS) 1 = PROGRAM/ERASE Attempted on a Locked Block; Operation Aborted 0 = No Operation to Locked Blocks	If a PROGRAM or ERASE operation is attempted to one of the locked blocks, this is set by the WSM. The operation specified is aborted and the device is returned to read status mode.
SR0	FAST PROGRAMMING ALGORITHM STATUS (FPAS) 1 = FPA PROGRAM/ERASE Busy 0 = FPA Ready	When this bit is set to "1," the FPA algorithm is active. When the FPA operation is complete, this bit is reset to "0."



**Figure 4**  
**Automated Word Programming**  
**Flowchart**



**FULL STATUS REGISTER CHECK FLOW**



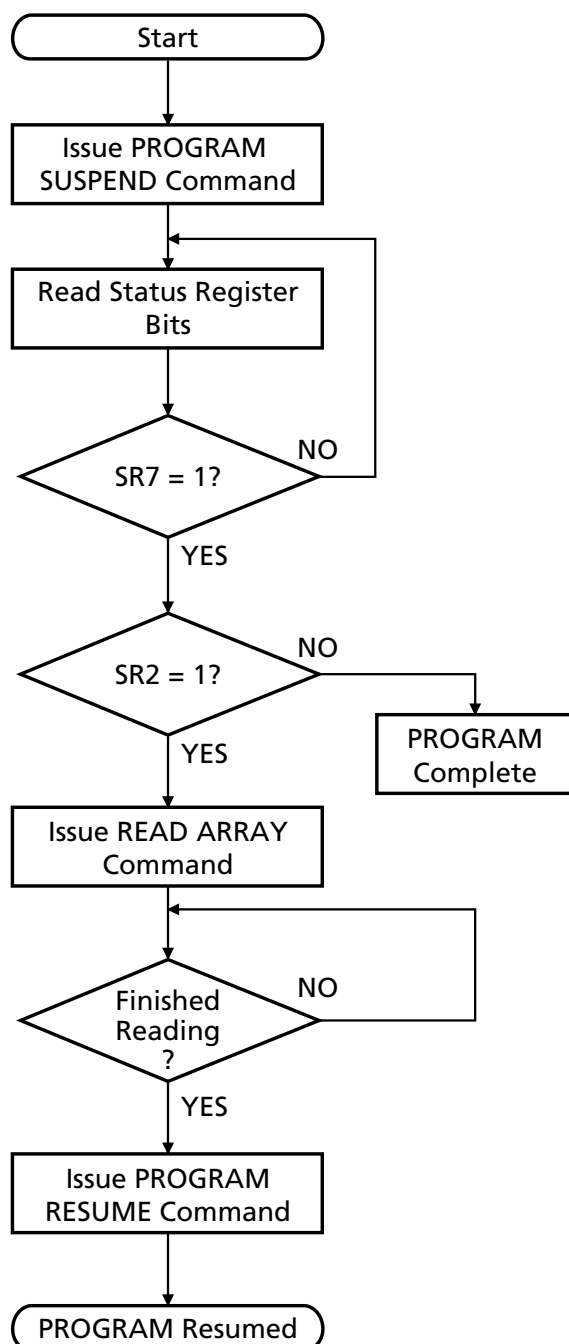
BUS OPERATION	COMMAND	COMMENTS
WRITE	WRITE PROGRAM SETUP	Data = 40h or 10h Addr = Address of word to be programmed
WRITE	WRITE DATA	Data = Word to be programmed Addr = Address of word to be programmed
READ		Status register data Toggle OE# or CE# to update status register.
Standby		Check SR7 1 = Ready, 0 = Busy
Repeat for subsequent words. Write FFh after the last word programming operation to reset the device to read array mode.		

BUS OPERATION	COMMAND	COMMENTS
Standby		Check SR1 1 = Detect locked block
Standby		Check SR3 <sup>2</sup> 1 = Detect V <sub>PP</sub> LOW
Standby		Check SR4 <sup>3</sup> 1 = Word program error

- NOTE:**
1. Full status register check can be done after each word or after a sequence of words.
  2. SR3 must be cleared before attempting additional PROGRAM/ERASE operations.
  3. SR4 is cleared only by the CLEAR STATUS REGISTER command, but it does not prevent additional program operation attempts.



**Figure 5**  
**PROGRAM SUSPEND/  
PROGRAM RESUME Flowchart**

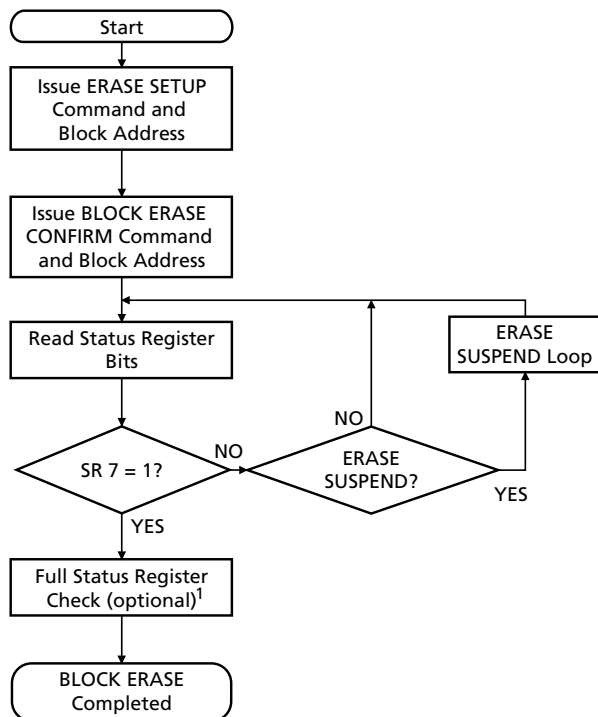


BUS OPERATION	COMMAND	COMMENTS
WRITE	PROGRAM SUSPEND	Data = B0h
READ		Status register data Toggle OE# or CE# to update status register.
Standby		Check SR7 1 = Ready
Standby		Check SR2 1 = Suspended
WRITE	READ ARRAY	Data = FFh
READ		Read data from block other than that being programmed.
WRITE	PROGRAM RESUME	Data = D0h

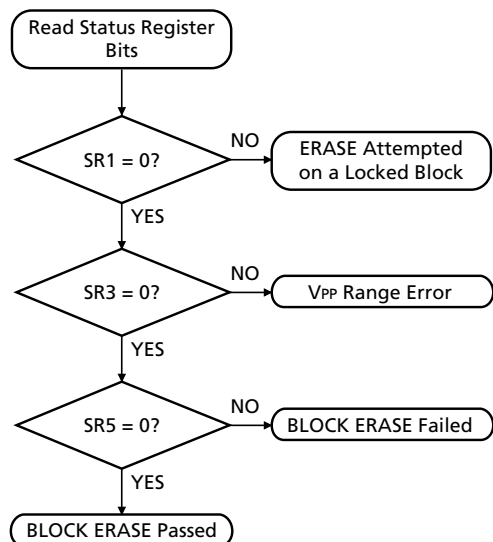
- NOTE:**
1. Full status register check can be done after each word or after a sequence of words.
  2. SR3 must be cleared before attempting additional PROGRAM/ERASE operations.
  3. SR5 is cleared only by the CLEAR STATUS REGISTER command in cases where multiple blocks are erased before full status is checked.



**Figure 6  
BLOCK ERASE Flowchart**



**FULL STATUS REGISTER CHECK FLOW**



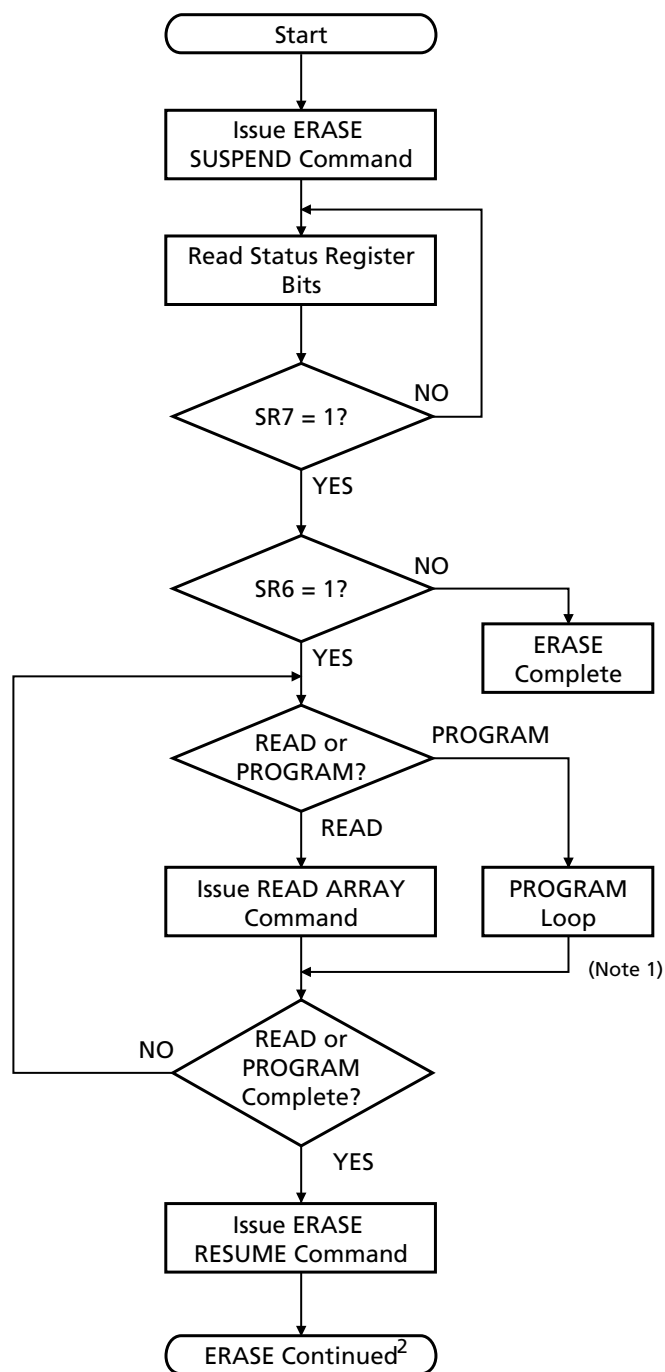
BUS OPERATION	COMMAND	COMMENTS
WRITE	WRITE ERASE SETUP	Data = 20h Block Addr = Address within block to be erased
WRITE	ERASE	Data = D0h Block Addr = Address within block to be erased
READ		Status register data Toggle OE# or CE# to update status register.
Standby		Check SR7 1 = Ready, 0 = Busy
Repeat for subsequent blocks. Write FFh after the last BLOCK ERASE operation to reset the device to read array mode.		

BUS OPERATION	COMMAND	COMMENTS
Standby		Check SR1 1 = Detect locked block
Standby		Check SR3 <sup>2</sup> 1 = Detect VPP block
Standby		Check SR5 <sup>3</sup> 1 = BLOCK ERASE error

- NOTE:**
1. Full status register check can be done after each block or after a sequence of blocks.
  2. SR3 must be cleared before attempting additional PROGRAM/ERASE operations.
  3. SR5 is cleared only by the CLEAR STATUS REGISTER command in cases where multiple blocks are erased before full status is checked.



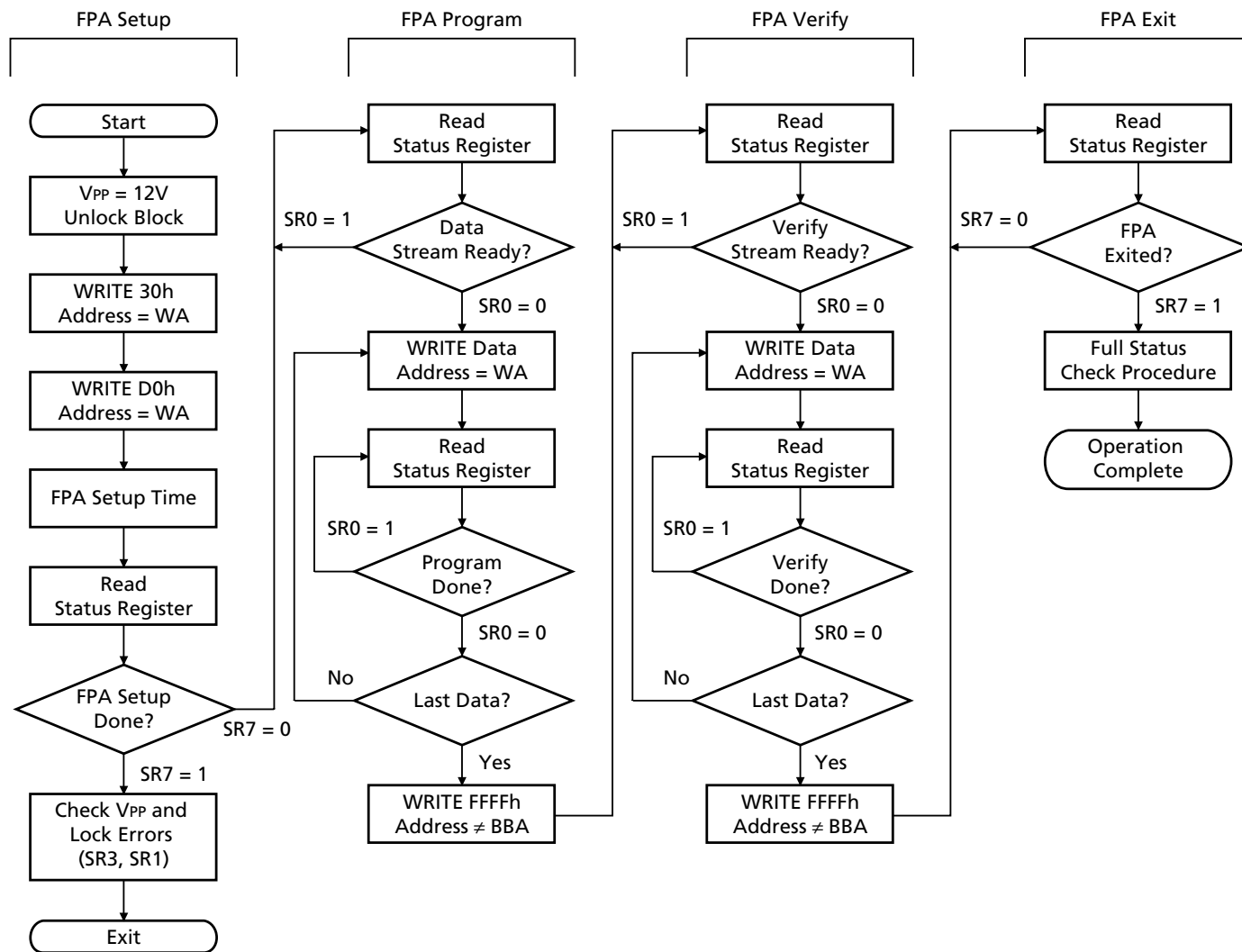
**Figure 7**  
**ERASE SUSPEND/ERASE RESUME**  
**Flowchart**



BUS OPERATION	COMMAND	COMMENTS
WRITE	ERASE SUSPEND	Data = 80h
READ		Status register data Toggle OE# or CE# to update status register.
Standby		Check SR7 1 = Ready
Standby		Check SR6 1 = Suspended
WRITE	READ ARRAY	Data = FFh
READ		Read data from block other than that being erased.
WRITE	ERASE RESUME	Data = D0h

**NOTE:** 1. See BLOCK ERASE Flowchart for complete erasure procedure.  
2. See Word Programming Flowchart for complete programming procedure.

**Figure 8**  
**Fast Programming Algorithm Flowchart**

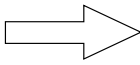
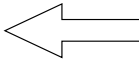


## READ-WHILE-WRITE/ERASE CONCURRENCY

It is possible for the device to read from one bank while erasing/writing to another bank. Once a bank enters the WRITE/ERASE operation, the other bank automatically enters read array mode. For example, during a READ CONCURRENCY operation, if a PROGRAM/ERASE command is issued in bank *a*, then bank *a* changes to the read status mode and bank *b* defaults to the read array mode. The device will read from bank *b* if the latched address resides in bank *b* (see Figure 9). Similarly, if a PROGRAM/ERASE command is issued in bank *b*, then bank *b* changes to read status mode and bank *a* defaults to read array mode. When returning to bank *a*, the device will read PROGRAM/ERASE status if the latched address resides in bank *a*. A correct bank address must be specified to read status register after returning from concurrent read in the other bank.

When reading the CFI or the chip protection register, concurrent operation is not allowed on the top boot device. Concurrent READ of the CFI or the chip protection register is only allowed when a PROGRAM or ERASE operation is performed on bank *b* on the bottom boot device. For a bottom boot device, reading of the CFI table or the chip protection register is only allowed if bank *b* is in read array mode. For a top boot device, reading of the CFI table or the chip protection register is only allowed if bank *a* is in read array mode.

**Figure 9**  
**READ-While-WRITE Concurrency**

Bank <i>a</i>	Bank <i>b</i>
1 - Erasing/writing to bank <i>a</i> 2 - Erasing in bank <i>a</i> can be suspended, and a WRITE to another block in bank <i>a</i> can be initiated. 3 - After the WRITE in that block is complete, an ERASE can be resumed by writing an ERASE RESUME command.	 1 - Reading from bank <i>b</i>
1 - Reading bank <i>a</i> 	1 - Erasing/writing to bank <i>b</i> 2 - Erasing in bank <i>b</i> can be suspended, and a WRITE to another block in bank <i>b</i> can be initiated. 3 - After the WRITE in that block is complete, an ERASE can be resumed by writing an ERASE RESUME command.

## BLOCK LOCKING

The MT28F322P3 Flash memory provides a flexible locking scheme which allows each block to be individually locked or unlocked with no latency.

The device offers two-level protection for the blocks. The first level allows software-only control of block locking (for data which needs to be changed frequently), while the second level requires hardware interaction before locking can be changed (code which does not require frequent updates).

Control signals WP#, DQ0, and DQ1 define the state of a block; for example, state [001] means WP# = 0, DQ0 = 0 and DQ1 = 1.

Table 8 defines all of the possible locking states.

**NOTE:** All blocks are software-locked upon power-up sequence completion.

### LOCKED STATE

After a power-up sequence completion, or after a reset sequence, all blocks are locked (states [001] or [101]). This means full protection from alteration. Any PROGRAM or ERASE operations attempted on a locked block will return an error on bit SR1 of the status register. The status of a locked block can be changed to unlocked or lock down using the appropriate software commands. Writing the lock command sequence, 60h followed by 01h, can lock an unlocked block.

### UNLOCKED STATE

Unlocked blocks (states [000], [100], [110]) can be programmed or erased. All unlocked blocks return to ERASE operations, but their protection status cannot be changed using software commands alone. A locked or unlocked block can be locked down by writing the lock down command sequence, 60h followed by 2Fh.

Locked down blocks revert to the locked state when the device is reset or powered down (see Table 4).

### LOCKED DOWN STATE

Blocks that are locked down (state [011]) are protected from PROGRAM and ERASE operations, but their protection status cannot be changed using software commands alone. A locked or unlocked block can be locked down by writing the lock down command sequence, 60h followed by 2Fh. Locked down blocks revert to the locked state when the device is reset or powered down.

The LOCK DOWN function is dependent on the WP# input. When WP# = 0, blocks in lock down [011] are protected from program, erase, and lock status changes. When WP# = 1, the lock down function is disabled ([111]), and locked down blocks can be individually unlocked by a software command to the [110] state, where they can be erased and programmed. These blocks can then be relocked [111] and unlocked [110] as desired while WP# remains HIGH. When WP# goes LOW, blocks that were previously locked down return to the locked down state [011] regardless of any changes made while WP# was HIGH. Device reset or power-down resets all locks, including those in lock down, to locked state (see Table 9).

### READING A BLOCK'S LOCK STATUS

The lock status of every block can be read in the read device identification mode. To enter this mode, write 90h to the bank containing address 00h. Subsequent READs at block address +00002 will output the lock status of that block. The lowest two outputs, DQ0 and DQ1, represent the lock status. DQ0 indicates the block lock/unlock status and is set by the LOCK com-

**Table 8  
Block Locking State Transition**

WP#	DQ1	DQ0	NAME	ERASE/PROGRAM ALLOWED	LOCK	UNLOCK	LOCK DOWN
0	0	0	Unlocked	Yes	To [001]	No Change	To [011]
0	0	1	Locked (Default)	No	No Change	To [000]	To [011]
0	1	1	Lock Down	No	No Change	No Change	No Change
1	0	0	Unlocked	Yes	To [101]	No Change	To [111]
1	0	1	Locked	No	No Change	To [100]	To [111]
1	1	0	Lock Down Disabled	Yes	To [111]	No Change	To [111]
1	1	1	Lock Down Disabled	No	No Change	To [110]	No Change

mand and cleared by the UNLOCK command. It is also automatically set when entering lock down. DQ1 indicates lock down status and is set by the LOCK DOWN command. It can only be cleared by reset or power-down, not by software. Table 8 shows the locking state transition scheme. The read array command, FFh, must be issued to the bank containing address 00h prior to issuing other commands.

### LOCKING OPERATIONS DURING ERASE SUSPEND

Changes to block lock status can be performed during an ERASE SUSPEND by using the standard locking command sequences to unlock, lock, or lock down. This is useful in the case when another block needs to be updated while an ERASE operation is in progress.

To change block locking during an ERASE operation, first write the ERASE SUSPEND command (B0h), then check the status register until it indicates that the ERASE operation has been suspended. Next, write the desired lock command sequence to block lock, and the lock status will be changed. After completing any desired LOCK, READ, or PROGRAM operations, resume the ERASE operation with the ERASE RESUME command (D0h).

If a block is locked or locked down during an ERASE SUSPEND on the same block, the locking status bits will be changed immediately. When the ERASE is resumed, the ERASE operation will complete.

A locking operation cannot be performed during a PROGRAM SUSPEND.

### CHIP PROTECTION REGISTER

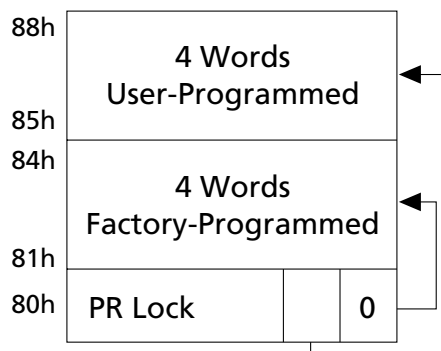
A 128-bit chip protection register can be used to fulfill the security considerations in the system (preventing the device substitution).

The 128-bit security area is divided into two 64-bit segments. The first 64 bits are programmed at the manufacturing site with a unique 64-bit unchangeable number. The other segment is left blank for customers to program as desired. (See Figure 10).

### READING THE CHIP PROTECTION REGISTER

The chip protection register is read in the device identification mode. To enter this mode, load the 90h command to the bank containing address 00h. Once in this mode, READ cycles from addresses shown in Table 9 retrieve the specified information. To return to the

**Figure 10  
Protection Register Memory Map**



**Table 9  
Chip Configuration Addressing<sup>1</sup>**

ITEM	ADDRESS <sup>2</sup>	DATA
Manufacturer Code (x16)	00000h	002Ch
Device Code <ul style="list-style-type: none"> <li>• Top boot configuration</li> <li>• Bottom boot configuration</li> </ul>	00001h	4494h 4495h
Block Lock Configuration <ul style="list-style-type: none"> <li>• Block is unlocked</li> <li>• Block is locked</li> <li>• Block is locked down</li> </ul>	XX002h	Lock DQ0 = 0 DQ0 = 1 DQ1 = 1
Chip Protection Register Lock	80h	PR Lock
Chip Protection Register 1	81h–84h	Factory Data
Chip Protection Register 2	85h–88h	User Data

**NOTE:** 1. Other locations within the configuration address space are reserved by Micron for future use.  
2. "XX" specifies the block address of lock configuration.

read array mode, write the READ ARRAY command (FFh). The read array command, FFh, must be issued to the bank containing address 00h prior to issuing other commands.

### ASYNCHRONOUS READ CYCLE

When accessing addresses in a random order or when switching between pages, the access time is given by  $t_{AA}$ .

When CE# and OE# are LOW, the data is placed on the data bus and the processor can read the data.

### PAGE READ MODE

The initial portion of the page mode cycle is the same as the asynchronous access cycle. Holding CE# LOW and toggling addresses A0–A2 allows random access of other words in the page. The page word size is eight words.

### V<sub>PP</sub> / V<sub>CC</sub> PROGRAM AND ERASE VOLTAGES

The MT28F322P3 Flash memory provides in-system programming and erase with V<sub>PP</sub> in the 1.8V–3.3V range. The 12V V<sub>PP</sub> mode programming is offered for compatibility with existing programming equipment, but does not enhance programming performance using the standard programming commands.

The device can withstand 100,000 WRITE/ERASE operations when V<sub>PP</sub> = V<sub>CC</sub> or 100 WRITE/ERASE operations and 10 cumulative hours when V<sub>PP</sub> = 12V.

In addition to the flexible block locking, the V<sub>PP</sub> programming voltage can be held low for absolute hardware write protection of all blocks in the flash device. When V<sub>PP</sub> is below V<sub>PLK</sub>, any PROGRAM or ERASE operation will result in an error, prompting the corresponding status register bit (SR3) to be set.

During WRITE and ERASE operations, the WSM monitors the V<sub>PP</sub> voltage level. WRITE/ERASE operations are allowed only when V<sub>PP</sub> is within the ranges specified in Table 10.

When V<sub>CC</sub> is below V<sub>LKO</sub>, any WRITE/ERASE operation will be disabled.

**Table 10**  
**V<sub>PP</sub> Range (V)**

	<b>MIN</b>	<b>MAX</b>
In-System	1.8	3.3
In-Factory	11.4	12.6

### STANDBY MODE

Icc supply current is reduced by applying a logic HIGH level on CE# and RST# to enter the standby mode. In the standby mode, the outputs are placed in High-Z. Applying a CMOS logic HIGH level on CE# and RST# reduces the current to Icc3 (MAX). If the device is deselected during an ERASE operation or during programming, the device continues to draw current until the operation is complete.

### AUTOMATIC POWER SAVE (APS) MODE

Substantial power savings are realized during periods when the array is not being read and the device is in the active mode. During this time the device switches to the automatic power save mode. When the device switches to this mode, Icc is reduced to a level comparable to Icc3. Further power savings can be realized by applying a logic HIGH level on CE# to place the device in standby mode. The low level of power is maintained until another operation is initiated. In this mode, the I/Os retain the data from the last memory address read until a new address is read. This mode is entered automatically if no address or control signals toggle.

### DEVICE RESET

To correctly reset the MT28F322P3 Flash memory, the RST# signal must be asserted (RST# = V<sub>IL</sub>) for a minimum of  $t_{RP}$ . After reset, the device can be accessed for a READ operation with a delayed access time of  $t_{RWH}$  from the rising edge of RST#. The circuitry used for generating the RST# signal needs to be common with the rest of the system reset to ensure that correct system initialization occurs. Please refer to the timing diagram for further details.

### POWER-UP SEQUENCE

The following power-up sequence is recommended to initialize internal chip operations:

- At power-up, RST# should be kept at V<sub>IL</sub> for 2μs after V<sub>CC</sub> reaches V<sub>CC</sub> (MIN).
- V<sub>CCQ</sub> should not come up before V<sub>CC</sub>.
- V<sub>PP</sub> should be kept at V<sub>IL</sub> to maximize data integrity.

When the power-up sequence is completed, RST# should be brought to V<sub>IH</sub>. To ensure a proper power-up, the rise time of RST# (10%–90%) should be <10μs.

**ABSOLUTE MAXIMUM RATINGS\***

Voltage to Any Ball Except V<sub>CC</sub> and V<sub>PP</sub>  
with Respect to V<sub>SS</sub> ..... -0.5V to +4V  
V<sub>PP</sub> Voltage (for BLOCK ERASE and PROGRAM  
with Respect to V<sub>SS</sub>) ..... -0.5V to +13.5V\*\*  
V<sub>CC</sub> and V<sub>CCQ</sub> Supply Voltage  
with Respect to V<sub>SS</sub> ..... -0.3V to +4V  
Output Short Circuit Current ..... 100mA  
Operating Temperature Range ..... -40°C to +85°C  
Storage Temperature Range ..... -55°C to +125°C  
Soldering Cycle ..... 260°C for 10s

\*Stresses greater than those listed under “Absolute Maximum Ratings” may cause permanent damage to the device. This is a stress rating only and functional operation of the device at these or any other conditions above those indicated in the operational sections of this specification is not implied. Exposure to absolute maximum rating conditions for extended periods may affect reliability.

\*\*Maximum DC voltage on V<sub>PP</sub> may overshoot to +13.5V for periods <20ns.

**RECOMMENDED OPERATING CONDITIONS**

PARAMETER		SYMBOL	MIN	MAX	UNITS	NOTE
Operating temperature		t <sub>A</sub>	-40	+85	°C	
V <sub>CC</sub> supply voltage		V <sub>CC</sub>	2.7	3.3	V	
I/O supply voltage		V <sub>CCQ</sub>	2.2	3.3	V	
V <sub>PP</sub> voltage		V <sub>PP1</sub>	1.8	3.3	V	
V <sub>PP</sub> in-factory programming voltage		V <sub>PP2</sub>	11.4	12.6	V	
Block erase cycling	V <sub>PP</sub> = V <sub>PP1</sub>	V <sub>PP1</sub>	–	100,000	Cycles	
	V <sub>PP</sub> = V <sub>PP2</sub>	V <sub>PP2</sub>	–	100	Cycles	1

**NOTE:** 1. V<sub>PP</sub> = V<sub>PP2</sub> is a maximum of 10 cumulative hours.



## DC CHARACTERISTICS<sup>1</sup>

PARAMETER	SYMBOL	MIN	TYP	MAX	UNITS	NOTES
Input Low Voltage	V <sub>IL</sub>	0	–	0.4	V	2
Input High Voltage	V <sub>IH</sub>	V <sub>CCQ</sub> - 0.4V	–	V <sub>CCQ</sub>	V	2
Output Low Voltage I <sub>OL</sub> = 100μA	V <sub>OL</sub>	-0.10	–	0.10	V	
Output High Voltage I <sub>OH</sub> = -100μA	V <sub>OH</sub>	V <sub>CCQ</sub> - 0.1V	–	–	V	
V <sub>PP</sub> Lockout Voltage	V <sub>PPLK</sub>	–	–	1	V	
V <sub>PP</sub> During PROGRAM/ERASE Operations	V <sub>PP1</sub>	1.8	–	3.3	V	
	V <sub>PP2</sub>	11.4	–	12.6	V	
V <sub>CC</sub> Program/Erase Lock Voltage	V <sub>LKO</sub>	1	–	–	V	
Input Leakage Current	I <sub>L</sub>	–	–	1	μA	
Output Leakage Current	I <sub>OZ</sub>	0.2	–	1	μA	
V <sub>CC</sub> Read Current Asynchronous Random Read, 70ns cycle Asynchronous Random Read, 200ns cycle	I <sub>CC1</sub>	– –	– –	15 6	mA mA	3, 4
V <sub>CC</sub> Page Mode Read Current at 70ns/30ns	I <sub>CC2</sub>	–	–	7	mA	3, 4
V <sub>CC</sub> Standby Current	I <sub>CC3</sub>	–	15μA	50	μA	
V <sub>CC</sub> Program Current	I <sub>CC4</sub>	–	10mA	20	mA	
V <sub>CC</sub> Erase Current	I <sub>CC5</sub>	–	15mA	25	mA	
V <sub>CC</sub> Erase Suspend Current	I <sub>CC6</sub>	–	15μA	50	μA	5
V <sub>CC</sub> Program Suspend Current	I <sub>CC7</sub>	–	15μA	50	μA	5
Read-While-Write Current	I <sub>CC8</sub>	–	–	40	mA	
V <sub>PP</sub> Current (Read, Standby, Erase Suspend, Program Suspend) V <sub>PP</sub> = V <sub>PP1</sub> V <sub>PP</sub> = V <sub>PP2</sub>	I <sub>PP1</sub>	– –	– –	1 200	μA μA	

- NOTE:**
1. All currents are in RMS unless otherwise noted.
  2. V<sub>IL</sub> may decrease to -0.4V, and V<sub>IH</sub> may increase to V<sub>CCQ</sub> + 0.3V for durations not to exceed 20ns.
  3. APS mode reduces I<sub>CC</sub> to approximately I<sub>CC3</sub> levels.
  4. Test conditions: V<sub>CC</sub> = V<sub>CC</sub> (MAX), CE# = V<sub>IL</sub>, OE# = V<sub>IH</sub>. All other inputs = V<sub>IH</sub> or V<sub>IL</sub>.
  5. I<sub>CC6</sub> and I<sub>CC7</sub> values are valid when the device is deselected. Any READ operation performed while in suspend mode will have an additional current draw of suspend current (I<sub>CC6</sub> or I<sub>CC7</sub>).



## CAPACITANCE

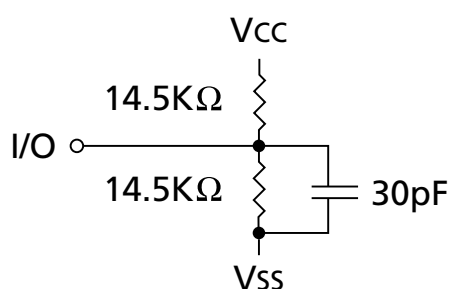
( $T_A = +25^\circ\text{C}$ ;  $f = 1\text{ MHz}$ )

PARAMETER/CONDITION	SYMBOL	TYP	MAX	UNITS
Input Capacitance	C	7	12	pF
Output Capacitance	C <sub>OUT</sub>	9	12	pF

## READ CYCLE TIMING REQUIREMENTS

PARAMETER	SYMBOL	-70		-80		UNITS
		MIN	MAX	MIN	MAX	
Address to output delay	$t_{AA}$		70		80	ns
CE# LOW to output delay	$t_{ACE}$		70		80	ns
Page address access	$t_{APA}$		30		30	ns
OE# LOW to output delay	$t_{AOE}$		25		25	ns
RST# HIGH to output delay	$t_{RWH}$		200		200	ns
RST# LOW pulse width	$t_{RP}$	100		100		ns
CE# or OE# HIGH to output High-Z	$t_{OD}$		25		25	ns
Output hold from address, CE# or OE# change	$t_{OH}$	0		0		ns
READ cycle time	$t_{RC}$		70		80	ns

**Figure 11  
Output Load Circuit**

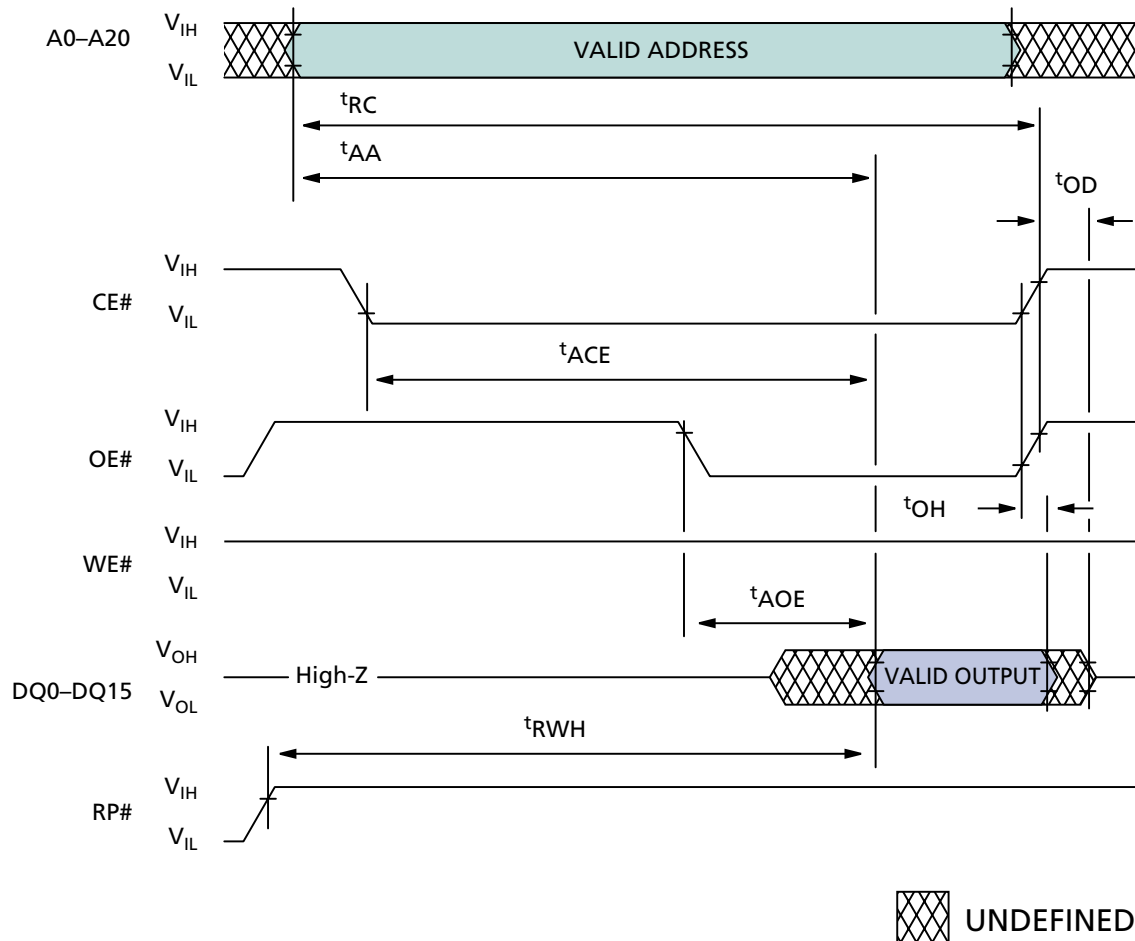


**WRITE CYCLE TIMING REQUIREMENTS**

PARAMETER	SYMBOL	-70/-80		UNITS
		MIN	MAX	
t <sub>RST#</sub> HIGH recovery to WE# going LOW	t <sub>RS</sub>	150		ns
CE# setup to WE# going LOW	t <sub>CS</sub>	0		ns
Write pulse width	t <sub>WP</sub>	30		ns
Data setup to WE# going HIGH	t <sub>DS</sub>	30		ns
Address setup to WE# going HIGH	t <sub>AS</sub>	30		ns
CE# hold from WE# HIGH	t <sub>CH</sub>	0		ns
Data hold from WE# HIGH	t <sub>DH</sub>	0		ns
Address hold from WE# HIGH	t <sub>AH</sub>	0		ns
Write pulse width HIGH	t <sub>WPH</sub>	30		ns
WP# setup to WE# going HIGH	t <sub>RHS</sub>	0		ns
V <sub>PP</sub> setup to WE# going HIGH	t <sub>VPS</sub>	200		ns
Write recovery before READ	t <sub>WOS</sub>	50		ns
WP# hold from valid SRD	t <sub>RHH</sub>	0		ns
V <sub>PP</sub> hold from valid SRD	t <sub>VPH</sub>	0		ns
WE# HIGH to data valid	t <sub>WB</sub>		t <sub>AA</sub> + 50	ns

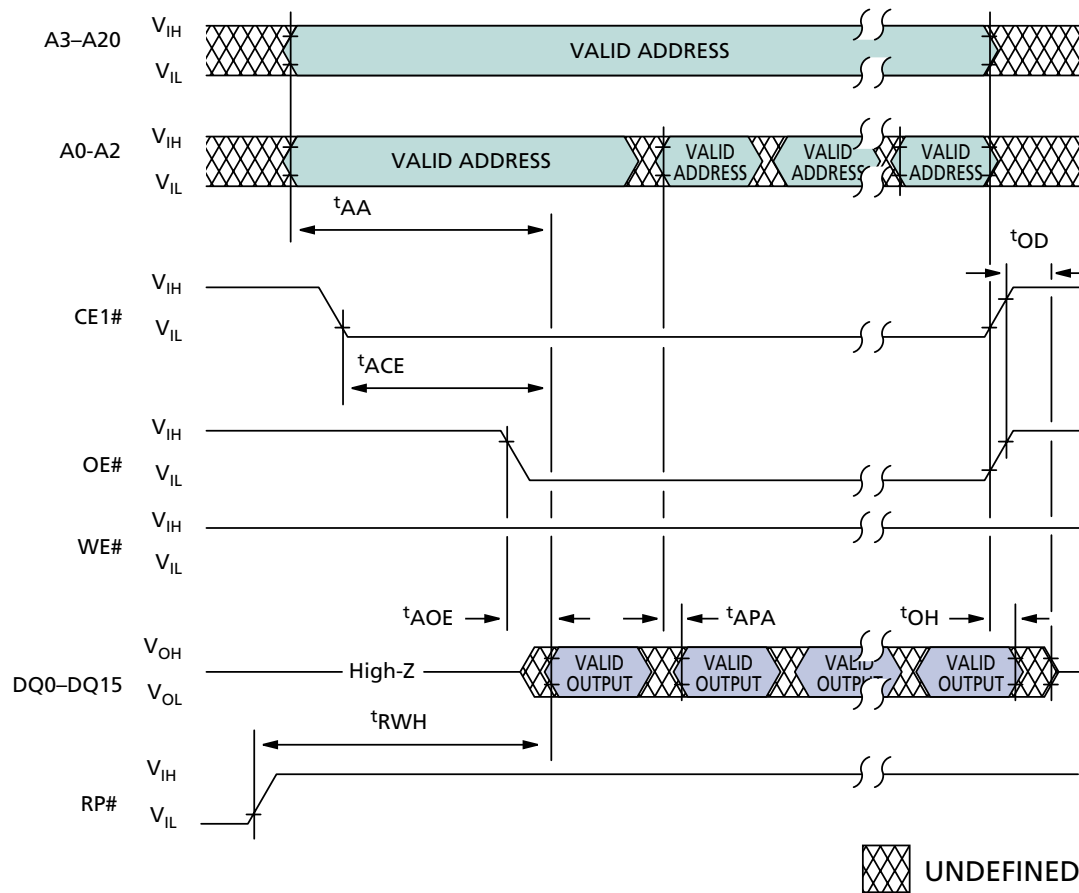
**ERASE AND PROGRAM CYCLE TIMING REQUIREMENTS**

PARAMETER	-70/-80		UNITS
	TYP	MAX	
4KW parameter block program time	40	800	ms
32KW parameter block program time	320	6,400	ms
Word program time	8	10,000	μs
4KW parameter block erase time	0.3	6	s
32KW parameter block erase time	0.5	6	s
Program suspend latency	5	10	μs
Erase suspend latency	5	20	μs
Chip programming time (FPA)		20	s
FPA setup time	5		μs

**SINGLE ASYNCHRONOUS READ OPERATION**

**READ TIMING PARAMETERS**

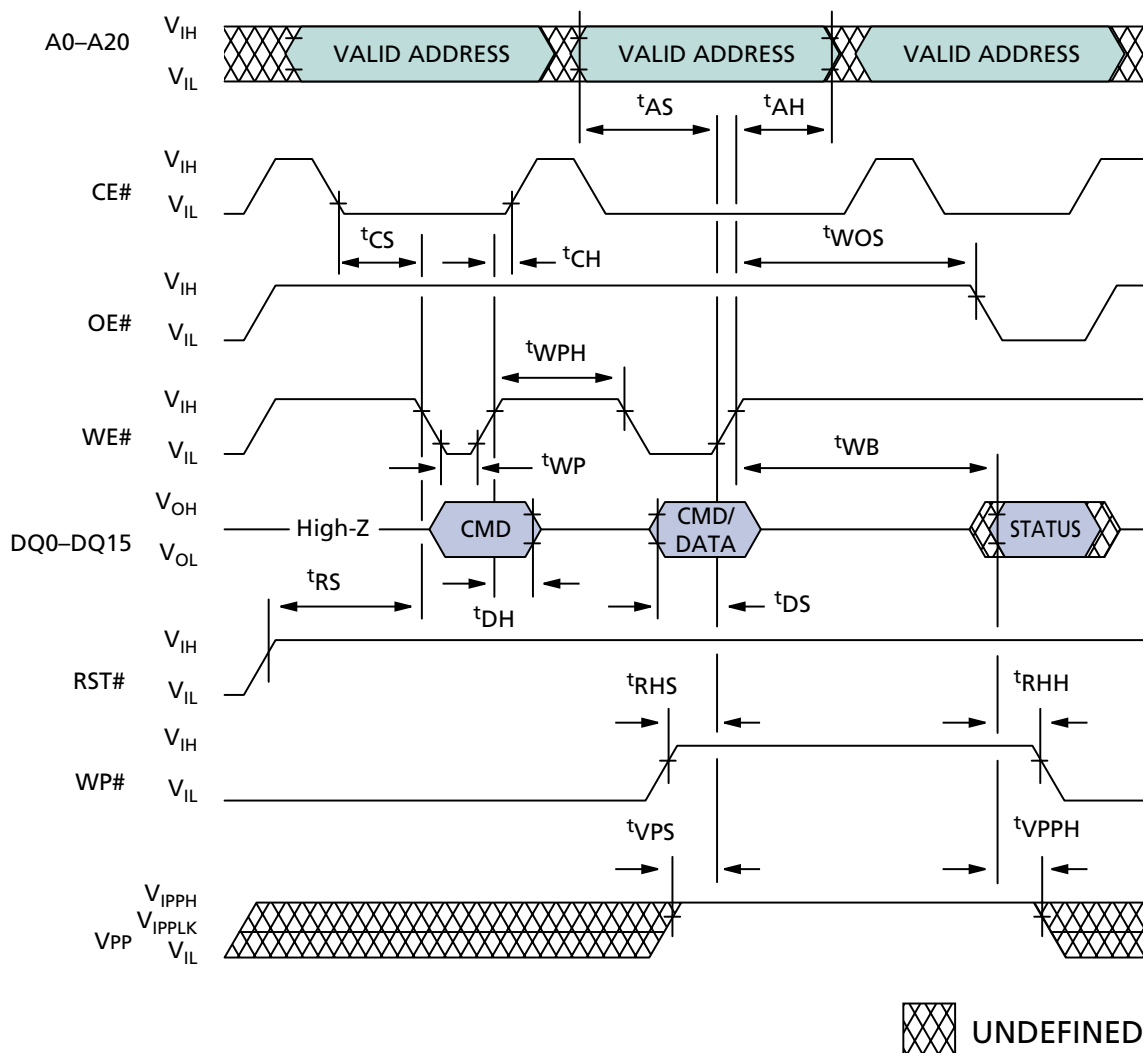
SYMBOL	-70		-80		UNITS
	MIN	MAX	MIN	MAX	
$t_{AA}$		70		80	ns
$t_{ACE}$		70		80	ns
$t_{AOE}$		25		25	ns
$t_{RWH}$		200		200	ns

SYMBOL	-70		-80		UNITS
	MIN	MAX	MIN	MAX	
$t_{OD}$		25		25	ns
$t_{OH}$	0		0		ns
$t_{RC}$		70		80	ns

**ASYNCHRONOUS PAGE MODE READ OPERATION**

**READ TIMING PARAMETERS**

SYMBOL	-70		-80		UNITS
	MIN	MAX	MIN	MAX	
$t_{AA}$		70		80	ns
$t_{ACE}$		70		80	ns
$t_{APA}$		30		30	ns
$t_{AOE}$		25		25	ns

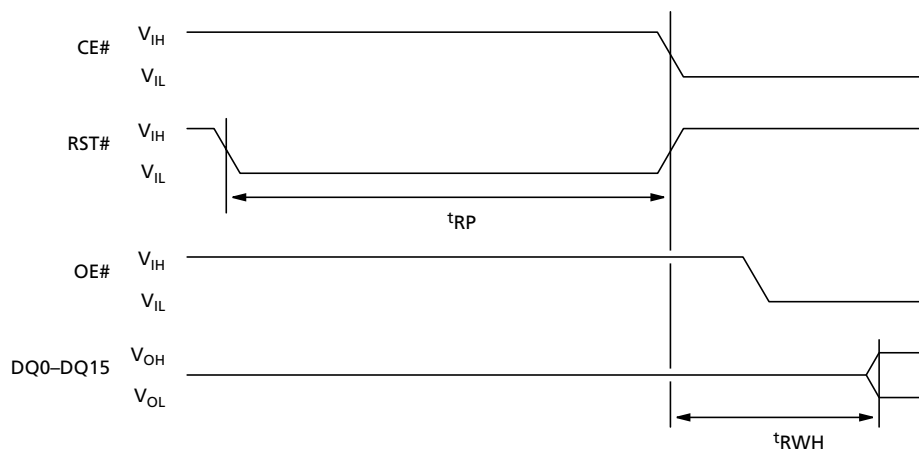
SYMBOL	-70		-80		UNITS
	MIN	MAX	MIN	MAX	
$t_{RWH}$		200		200	ns
$t_{OD}$		25		25	ns
$t_{OH}$	0		0		ns

**TWO-CYCLE PROGRAMMING/ERASE OPERATION**

**WRITE TIMING PARAMETERS**

SYMBOL	-70/-80		UNITS
	MIN	MAX	
$t_{RS}$	150		ns
$t_{CS}$	0		ns
$t_{WP}$	30		ns
$t_{DS}$	30		ns
$t_{AS}$	30		ns
$t_{CH}$	0		ns
$t_{DH}$	0		ns

SYMBOL	-70/-80		UNITS
	MIN	MAX	
$t_{AH}$	0		ns
$t_{RHS}$	0		ns
$t_{VPS}$	200		ns
$t_{WOS}$	50		ns
$t_{RHH}$	0		ns
$t_{VPPH}$	0		ns
$t_{WB}$		$t_{AA} + 50$	ns

### RESET OPERATION



### READ AND WRITE TIMING PARAMETERS

SYMBOL	-70/-80		UNITS
	MIN	MAX	
$t_{RWH}$		200	ns
$t_{RP}$	100		ns


**Table 11  
CFI**

OFFSET	DATA	DESCRIPTION
00	2Ch	Manufacturer code
01	94h	Top boot block device code
	95h	Bottom boot block device code
02–0F	reserved	Reserved
10, 11	0051, 0052	"QR"
12	0059	"Y"
13, 14	0003, 0000	Primary OEM command set
15, 16	0039, 0000	Address for primary extended table
17, 18	0000, 0000	Alternate OEM command set
19, 1A	0000, 0000	Address for OEM extended table
1B	0027	V <sub>CC</sub> MIN for Erase/Write; Bit7–Bit4 Volts in BCD; Bit3–Bit0 100mV in BCD
1C	0033	V <sub>CC</sub> MAX for Erase/Write; Bit7–Bit4 Volts in BCD; Bit3–Bit0 100mV in BCD
1D	00B4	V <sub>PP</sub> MIN for Erase/Write; Bit7–Bit4 Volts in Hex; Bit3–Bit0 100mV in BCD
1E	00C6	V <sub>PP</sub> MAX for Erase/Write; Bit7–Bit4 Volts in Hex; Bit3–Bit0 100mV in BCD
1F	0003	Typical timeout for single byte/word program, 2 <sup>n</sup> μs, 0000 = not supported
20	0000	Typical timeout for maximum size multiple byte/word program, 2 <sup>n</sup> μs, 0000 = not supported
21	0009	Typical timeout for individual block erase, 2 <sup>n</sup> ms, 0000 = not supported
22	0000	Typical timeout for full chip erase, 2 <sup>n</sup> ms, 0000 = not supported
23	000C	Maximum timeout for single byte/word program, 2 <sup>n</sup> μs, 0000 = not supported
24	0000	Maximum timeout for maximum size multiple byte/word program, 2 <sup>n</sup> μs, 0000 = not supported
25	0003	Maximum timeout for individual block erase, 2 <sup>n</sup> ms, 0000 = not supported
26	0000	Maximum timeout for full chip erase, 2 <sup>n</sup> ms, 0000 = not supported
27	0016	Device size, 2 <sup>n</sup> bytes
28	0001	Bus interface x16 = 1
29	0000	Flash device interface description 0000 = async
2A, 2B	0000, 0000	Maximum number of bytes in multi-byte program or page, 2 <sup>n</sup>
2C	0003	Number of erase block regions within device (4K words and 32K words)
2D, 2E	002F, 0000	Top boot block device erase block region information 1, 8 blocks ...
	0007, 0000	Bottom boot block device erase block region information 1, 8 blocks ...
2F, 30	0000, 0001	Top boot block device.....of 8KB
	0020, 0000	Bottom boot block device.....of 8KB
31, 32	000E, 0000	Top boot block 15 blocks of ....
	000E, 0000	Bottom boot block 15 blocks of ....
33, 34	0000, 0001	.....64KB

(continued on the next page)



**Table 11  
CFI (continued)**

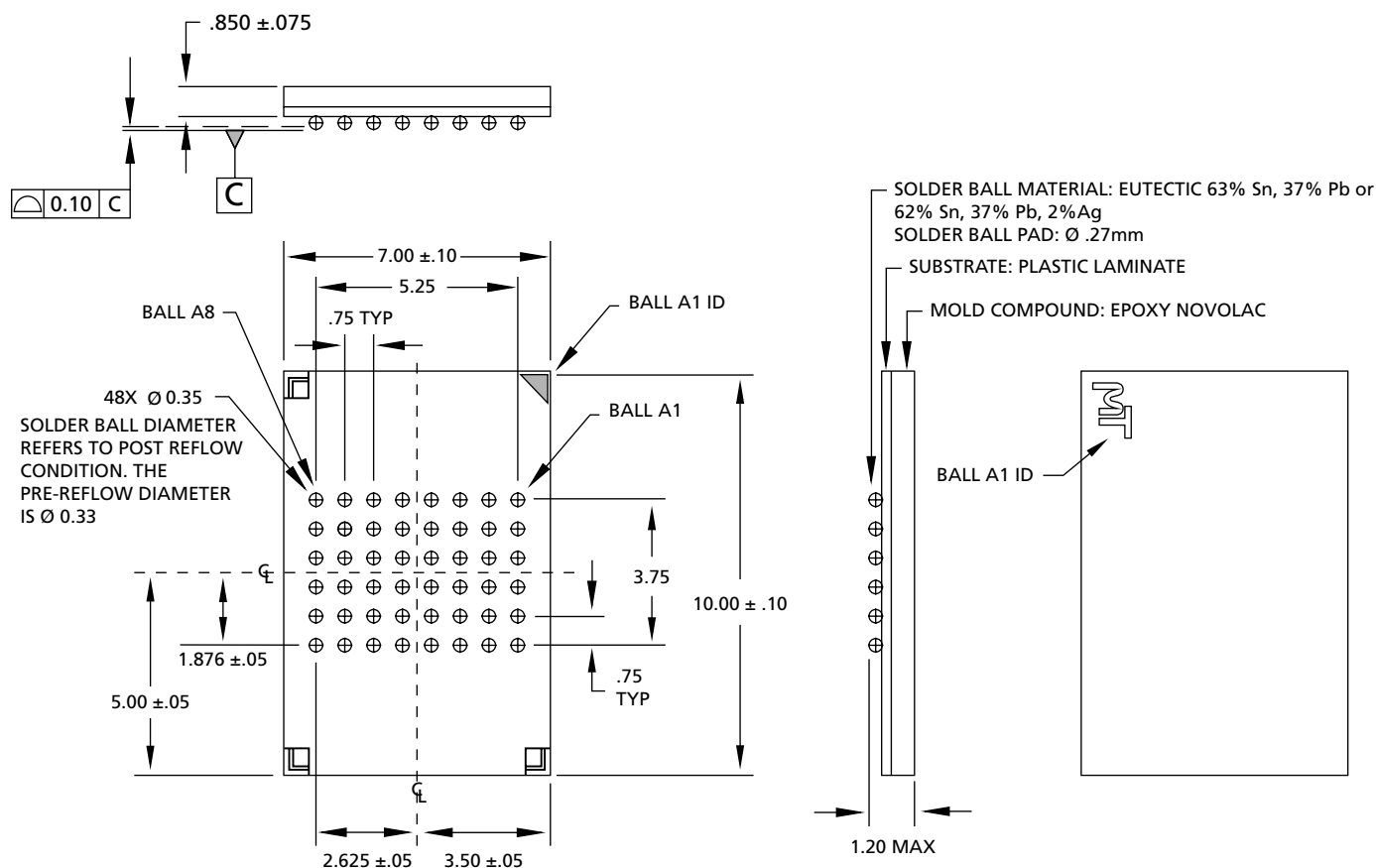
OFFSET	DATA	DESCRIPTION
35, 36	0007, 0000	Top boot block device.....48 blocks of
	002F, 0000	Bottom boot block device.....48 blocks of
37, 38	0020, 0000	Top boot block device.....64KB
	0000, 0001	Bottom boot block device.....64KB
39, 3A	0050, 0052	"PR"
3B	0049	"I"
3C	0030	Major version number, ASCII
3D	0031	Minor version number, ASCII
3E 3F 40 41	00E6 0002 0000 0000	Optional Feature and Command Support Bit 0 Chip erase supported no = 0 Bit 1 Suspend erase supported = yes = 1 Bit 2 Suspend program supported = yes = 1 Bit 3 Chip lock/unlock supported = no = 0 Bit 4 Queued erase supported = no = 0 Bit 5 Instant individual block locking supported = yes = 1 Bit 6 Protection bits supported = yes = 1 Bit 7 Page mode read supported = yes = 1 Bit 8 Synchronous read supported = no = 0 Bit 9 Simultaneous operation supported = yes = 1
42	0001	Program supported after erase suspend = yes
43, 44	0003, 0000	Bit 0 block lock status active = yes; Bit 1 block lock down active = yes
45	0030	V <sub>CC</sub> supply optimum, 00 = not supported, Bit7–Bit4 Volts in BCD; Bit3–Bit0 100mV in BCD
46	00C0	V <sub>PP</sub> supply optimum, 00 = not supported, Bit7–Bit4 Volts in BCD; Bit3–Bit0 100mV in BCD
47	0001	Number of protection register fields in JEDEC ID space
48, 49	0080, 0000	Lock bytes LOW address, lock bytes HIGH address
4A, 4B	0003, 0003	2 <sup>n</sup> factory programmed bytes, 2 <sup>n</sup> user programmable bytes
4C	0003	Background Operation 0000 = Not used 0001 = 4% block split 0002 = 12% block split 0003 = 25% block split 0004 = 50% block split
4D	0000	Burst Mode Type 0000 = No burst mode 00x1 = 4 words MAX 00x2 = 8 words MAX 00x3 = 16 words MAX 001x = Linear burst, and/or 002x = Interleaved burst, and/or 004x = Continuous burst

(continued on the next page)



**Table 11**  
**CFI (continued)**

OFFSET	DATA	DESCRIPTION
4E	0002	Page Mode Type 0000 = No page mode 0001 = 4-word page 0002 = 8-word page 0003 = 16-word page 0004 = 32-word page
4F	0000	Not used

**48-BALL FBGA**


**NOTE:** 1. All dimensions in millimeters.  
2. Package width and length do not include mold protrusion; allowable mold protrusion is 0.27mm per side.

**DATA SHEET DESIGNATION**

Preliminary: This data sheet contains initial characterization limits that are subject to change upon full characterization of production devices.



8000 S. Federal Way, P.O. Box 6, Boise, ID 83707-0006, Tel: 208-368-3900

E-mail: [prodmktg@micron.com](mailto:prodmktg@micron.com), Internet: <http://www.micron.com>, Customer Comment Line: 800-932-4992

Micron and the M logo are registered trademarks, and the Micron logo is a trademark of Micron Technology, Inc.



## REVISION HISTORY

Rev. 3, PRELIMINARY .....	7/02
• Updated Status Register section	
• Updated command descriptions	
• Updated flowcharts	
• Updated Read-While-Write/EraseConcurrence section	
• Updated timing diagrams	
Rev. 2, PRELIMINARY .....	4/02
• Changed <sup>t</sup> AH from 9ns to 0ns	
• Updated READ CHIP PROTECTION IDENTIFICATION DATA text	
Rev. 2, PRELIMINARY .....	3/02
• Updated DC Characteristics table	
• Updated Table 10	
• Updated Asynchronous READ Cycle Timing Requirements table	
• Updated WRITE Cycle Timing Requirements table	
• Updated timing diagrams and parameters	
Original document, PRELIMINARY .....	1/02